Case 1: D.R. vs F.S. Hypotheses

Case 2: Path Verification

Case 3: Alibi Verification

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LOGIC PROGRAMMING AND NON-MONOTONIC REASONING

DIGITAL FORENSICS EVIDENCE ANALYSIS: AN ANSWER SET PROGRAMMING APPROACH FOR GENERATING INVESTIGATION HYPOTHESES

Stefania Costantini Giovanni De Gasperis Raffaele Olivieri

LEXINGTON, KY (USA) SEPTEMBER 28, 2015

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Who am I?

Ph.D. Univ. L'Aquila



Off. in Chief of a Unit of RaCIS



Digital Forenser

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Who Am I ?	Research Scope	Case 1: D.R. vs F.S. Hypotheses	Case 2: Path Verification	Case 3: Alibi Verification
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2 Case 1: D.R. vs F.S. Hypotheses

3 Case 2: Path Verification



Case 3: Alibi Verification

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2 Case 1: D.R. vs F.S. Hypotheses





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Research Scope ●000 Case 1: D.R. vs F.S. Hypotheses

Case 2: Path Verification

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Research Scope

Digital Forensics



admissibility

The methods must conform to the dictates of the Law

Definition

is a branch of *Criminalistics* which deals with the

- identification
- acquisition
- preservation
- analysis
- presentation

of the information content of ("*Digital Evidence*") with procedures resistant to any complaints in both civil and criminal court.

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Research Scope



 Identification 	\rightarrow	operator	\rightarrow	ISO 27037:2012	h
Acquisition	\rightarrow	dispositivi e tools	\rightarrow	Container(DD,EWF,AFF,)	I
3 Preservation	\rightarrow	HASH algorithms	\rightarrow	Warranty(Integrity/Authenticity)	I
4 Analysis	\rightarrow	tool + analyst	\rightarrow	Recovery e/o filtering	I
5 Presentation	\rightarrow	responsabile	\rightarrow	Valutatation and decision	
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Case 1: D.R. vs F.S. Hypotheses

Case 2: Path Verification

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Research Scope





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Research Scope

Analysis of "Digital Evidence"

Often, different technicians analyzing the same case reach different conclusions, and this may determine different judge's decisions in court.

The analysis of *Digital Evidence* often concerns the examination of incomplete knowledge and or fragmented, and complex scenarios, and includes:

- time evolution
- causation
- uncertainty and doubts
- randomness
- existence of alternative scenarios

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Digital Forensics Evidence Analysis: An Answer Set Programming Approach for Generating Investigation Hypotheses

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Case 3: Alibi Verification

Research Scope

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Research Scope

Answer Set Programming

- various investigation cases are reducible to known optimization problems, for which ASP is particularly suitable;
- ease of reading and interpretation
- stable model semantics
- every answer set represents a possible problem solution

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Who	Am		Research	Scope

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Case 2: Path Verification

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Data Recovery & File Sharing Hypotheses



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Case 1: D.R. vs F.S. Hypotheses

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Data Recovery & File Sharing Hypotheses



Marriage Problem

knowing the preferences of a group of guys and girls the problem is to pair people so that everyone has a preferred partner to the best possible degree.

Reduction

Men list includes NAMEs of the files contained in INDX files, while Women list consists in the list of name of recovered files. Preferences are elicited from metadata.

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Case 1: D.R. vs F.S. Hypotheses

Case 2: Path Verification

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Case 2: Path Verification

Case 3: Alibi Verification

Data Recovery & File Sharing Hypotheses

ASP code

```
likes(nome1,file1).
likes(nome1,file2).
likes(nome2.file1).
likes(nome2,file3).
likes(nome2,file2).
likes(nome3,file3).
likes(nome3,file2).
```

```
bigamia(X,Y) := likes(X,Y), likes(X,Y1),
coppia(X,Y), coppia(X,Y1), Y!=Y1.
bigamia(X,Y) :- likes(X,Y), likes(X1,Y),
coppia(X1,Y), X!=X1.
coppia(X,Y) := likes(X,Y), not bigamia(X,Y).
```

Risultato

Answer: 1 Stable Model: coppia(nome3,file2) coppia(nome2,file3) coppia(nome1,file1) Answer: 2 Stable Model: coppia(nome3,file3) coppia(nome2,file1) coppia(nome1,file2) Answer: 3 Stable Model: coppia(nome3,file3) coppia(nome2,file2) coppia(nome1,file1)

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Who Am	۱I	Research	Scop

Case 3: Alibi Verification

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Case 2: Path Verification

Case 3: Alibi Verification

Path Verification

Sexual Abuse





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Case 1: D.R. vs F.S. Hypothese

Case 2: Path Verification

Case 3: Alibi Verification

Path Verification



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Case 1: D.R. vs F.S. Hypotheses

Case 2: Path Verification

Case 3: Alibi Verification

Path Verification

Hidato Matrix						
18				26		
19			27			
	14			23	31	
1			8	33		
		5				
		10		36	35	

Hidato Problem

On a sparse matrix, the goal of Hidato is to fill the array using consecutive numbers in cells adjacent horizontally, vertically or diagonally, creating an ideal path.

Hidato Reduction

The matrix represents the geographic area of interest, where each matrix element represents an area traversable in a unit of time and the value represent the time.

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Case 1: D.R. vs F.S. Hypotheses

Case 2: Path Verification

Case 3: Alibi Verification

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Case 3: Alibi Verification

Path Verification

ASP Code

#const n = 6. matrix(1, 1, 18). matrix(1, 5, 26). matrix(2, 1, 19). matrix(2, 4, 27). matrix(3, 2, 14). matrix(3, 5, 23). matrix(3, 6, 31). matrix(4, 1, 1). matrix(4, 4, 8). matrix(4, 5, 33). matrix(5, 3, 5). matrix(6, 3, 10). matrix(6, 5, 36). matrix(6, 6, 35). size(1..n). values(1..n*n). values2(1..n*n-1). diffs(-1;0;1).

- 1 x(Row, Col, Value): values(Value) 1 :- size(Row), size(Col). 1 x(Row, Col, Value): size(Row): size(Col) 1 :- values(Value). x(Row, Col, Value):matrix(Row, Col, Value).
- valid(Row, Col, Row2, Col2) :- diffs(A), diffs(B), Row2 = Row+A, Col2 = Col+B, Row2 >= 1, Col2 >= 1, Row2 <= size, Col2 <= size, size(Row), size(Col).</pre>
- :- x(Row, Col, Value+1), x(Row2, Col2, Value), not valid(Row, Col, Row2, Col2), values2(Value).

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Path Verification

Results

Answer Set: 1

 $\begin{array}{l} x(1,1,18) \ x(1,5,26) \ x(2,1,19) \ x(2,4,27) \ x(3,2,14) \ x(3,5,23) \ x(3,6,31) \\ x(4,1,1) \ x(4,4,8) \ x(4,5,33) \ x(5,3,5) \ x(6,3,10) \ x(6,5,36) \ x(6,6,35) \ x(5,1,2) \\ x(6,1,3) \ x(6,2,4) \ x(6,4,6) \ x(5,5,7) \ x(5,4,9) \ x(5,2,11) \ x(4,2,12) \ x(3,1,13) \\ x(4,3,15) \ x(3,3,16) \ x(2,3,21) \ x(3,4,22) \ x(2,6,24) \ x(1,6,25) \ x(1,3,28) \\ x(1,4,29) \ x(2,5,30) \ x(4,6,32) \ x(5,6,34) \ x(1,2,20) \ x(2,2,17) \end{array}$

Answer Set: 2

 $\begin{array}{l} x(1,1,18) \ x(1,5,26) \ x(2,1,19) \ x(2,4,27) \ x(3,2,14) \ x(3,5,23) \ x(3,6,31) \\ x(4,1,1) \ x(4,4,8) \ x(4,5,33) \ x(5,3,5) \ x(6,3,10) \ x(6,5,36) \ x(6,6,35) \ x(5,1,2) \\ x(6,1,3) \ x(6,2,4) \ x(6,4,6) \ x(5,5,7) \ x(5,4,9) \ x(5,2,11) \ x(4,3,12) \ x(3,3,13) \\ x(4,2,15) \ x(3,1,16) \ x(2,3,21) \ x(3,4,22) \ x(2,6,24) \ x(1,6,25) \ x(1,3,28) \\ x(1,4,29) \ x(2,5,30) \ x(4,6,32) \ x(5,6,34) \ x(1,2,20) \ x(2,2,17) \end{array}$

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Research Scope

Case 2: Path Verification 000000

Case 3: Alibi Verification

2

Path Verification

Answer Set 1						
18	20	28	29	26	25	
19	17	21	27	30	24	
13	14	16	22	23	31	
1	12	15	8	33	32	
2	11	5	9	7	34	
3	4	10	6	36	35	

Answer Set 2						
18	20	28	29	26	25	
19	17	21	27	30	24	
16	14	13	22	23	31	
1	15	12	8	33	32	
2	11	5	9	7	34	
3	4	10	6	36	35	

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Who Am I ?	Research Scope	Case 1: D.R. vs F.S. Hypotheses	Case 2: Path Verification

Case 3: Alibi Verification

2





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Case 2: Path Verification

Case 3: Alibi Verification

Alibi Verification

Murder



Suspect Arrested



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Case 3: Alibi Verification

Alibi Verification

Alibi:

During interrogation the suspect said:

- left his home (place X) at a certain time;
- reached his office (place Y) where he worked on the PC;
- left the Office to go to his friend (place Z) where entering discovered the body;
- called the Police immediatly.

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Case 3: Alibi Verification

Alibi Verification

To verify the alibi were made the following analysis:

- smartphone's memories of the suspect;
- PC seized in his Office;
- a system of video surveillance installed at a post office, near the place Z, which recorded images of a very crowded street, where the investigators retrive several image sequences where compare a subject with charatteristics compatible with the suspect.

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Case 3: Alibi Verification

Alibi Verification

Monkey & Banana Problem

The "Monkey & Banana" problem is a typical planning problem. A monkey is in a room with a chair and a banana tied to the ceiling. The monkey can not reach the banana, unless it is located on the chair. Determine the correct sequence of actions.

Monkey & Banana Reduction

Monkey	
	Raise Alarm
Init.ial Position Monkey	X
	Z
Walks	Walks
Move Chair	Motion to Z
Ascend	
	Unknown Action

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Case 3: Alibi Verification

Alibi Verification

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A monkey is in a room with a chair and a banana tied to the ceiling. The monkey can not reach the banana, unless it is located on the chair. Determine the correct sequence of actions.

Monkey & Banana Reduction

Monkey Banana Eats Banana Init.ial Position Monkey Initial Position Chair Below Banana	$\begin{array}{c} \rightarrow \\ \rightarrow \\ \rightarrow \\ \rightarrow \\ \rightarrow \\ \rightarrow \end{array}$	Suspect Body Raise Alarm X Y Z
Walks Move Chair	\rightarrow \rightarrow	Walks Motion to Z
Ascend Idle	${\rightarrow}$ \rightarrow	Open the Door Unknown Action

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Case 3: Alibi Verification

Alibi Verification

Running ASP code we get, among others, the following answer sets that contradict the investigation thesis :

Answer Set 1

Answer Set 2

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Conclusions

Objectives

- demonstrate the applicability of Logic Programming and non-Monotonic Reasoning to Digital Forensics:
- convince the parties involved in the trial of the limitations of the current analysis techniques;
- provide, in the long term, to police, prosecutors, lawyers, judges, investigators, intelligence agencies, criminologists, etc., with a decision support systems to help them in their activities.

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