

Automated inference of rules with exception from past legal cases using ASP

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- ❑ Introduction and Motivation
- ❑ Formalisation
- ❑ ASP Workflow
- ❑ Evaluation
- ❑ Summary and Future Work

- In legal reasoning, we use written rules to make judgements
- Cases are very important since they can sometimes reveal exceptional situations not considered in written rules
- And these cases might have exceptions revealed by future cases

The purpose of this research is to:

- Formalise case rules in terms of general rules/exceptions.
- Investigate how to infer case-rules from previously judged cases

Motivation



	Factors	Factors	Judgement
			-DepriveBuyerOfEntitlement (default)

Motivation



	Factors	Factors	Judgement
			¬DepriveBuyerOfEntitlement (default)
1	DeliveredOnTime (dot)		¬DepriveBuyerOfEntitlement

Motivation



	Factors	Factors	Judgement
			¬DepriveBuyerOfEntitlement (default)
1	DeliveredOnTime (dot)		¬DepriveBuyerOfEntitlement
2	DeliveredOnTime (dot) ItemsWereDamanged (ooo)		DepriveBuyerOfEntitlement

Motivation



	Factors		Factors	Judgement
				¬DepriveBuyerOfEntitlement (default)
1	DeliveredOnTime	(dot)		¬DepriveBuyerOfEntitlement
2	DeliveredOnTime ItemsWereDamaged	(dot) (ooo)		DepriveBuyerOfEntitlement
3	DeliveredOnTime ItemsWereDamaged	(dot) (ooo)	DamagesRepairable (rpl) BuyerFixedRepairTime (far) ItemsRepairedInTime (ria)	¬DepriveBuyerOfEntitlement

Motivation



	Factors		Factors	Judgement
				¬DepriveBuyerOfEntitlement (default)
1	DeliveredOnTime	(dot)		¬DepriveBuyerOfEntitlement
2	DeliveredOnTime ItemsWereDamaged	(dot) (ooo)		DepriveBuyerOfEntitlement
3	DeliveredOnTime ItemsWereDamaged	(dot) (ooo)	DamagesRepairable (rpl) BuyerFixedRepairTime (far) ItemsRepairedInTime (ria)	¬DepriveBuyerOfEntitlement
4	DeliveredOnTime ItemsWereDamaged	(dot) (ooo)	DamagesRepairable (rpl) BuyerFixedRepairTime (far)	?

Motivation



Should judgement of case 4 be
DepriveBuyerOfEntitlement
because case 2?
or
Should the judgement of case 4 be
¬DepriveBuyerOfEntitlement
because case 3?

	Factors			DepriveBuyerOfEntitlement
1	DeliveredOnTime	(dot)		DepriveBuyerOfEntitlement (default)
2	DeliveredOnTime ItemsWereDamaged	(dot) (ooo)		DepriveBuyerOfEntitlement
3	DeliveredOnTime ItemsWereDamaged	(dot) (ooo)	DamagesRepairable (rpl) BuyerFixedRepairTime (far) ItemsRepairedInTime (ria)	¬DepriveBuyerOfEntitlement
4	DeliveredOnTime ItemsWereDamaged	(dot) (ooo)	DamagesRepairable (rpl) BuyerFixedRepairTime (far)	?

- A **case** is a subset of a set of factors F .
- A **case with judgement** is $cj = \langle c, j \rangle$, where $j \in \{+, -\}$
 $case(cj) = c$ and $judgement(cj) = j$
- A **case base** CB is a set of cases with judgements

All casebases include $\langle \emptyset, j_0 \rangle$

- \emptyset is the **empty case** and
- j_0 is the **default judgement**, assumed in the absence of any factor

For all cases cj_1 and cj_2 in CB ,
if $case(cj_1) = case(cj_2)$ then $judgement(cj_1) = judgement(cj_2)$

Raw Attack Relation (RA)

The **raw attack** relation

$$RA = \{ \langle cj_1, cj_2 \rangle \mid case(cj_1) \supset case(cj_2) \text{ and } judgement(cj_1) \neq judgement(cj_2) \}$$

cj_1 **raw attacks** cj_2 is denoted as $cj_1 \rightarrow_r cj_2$

Example

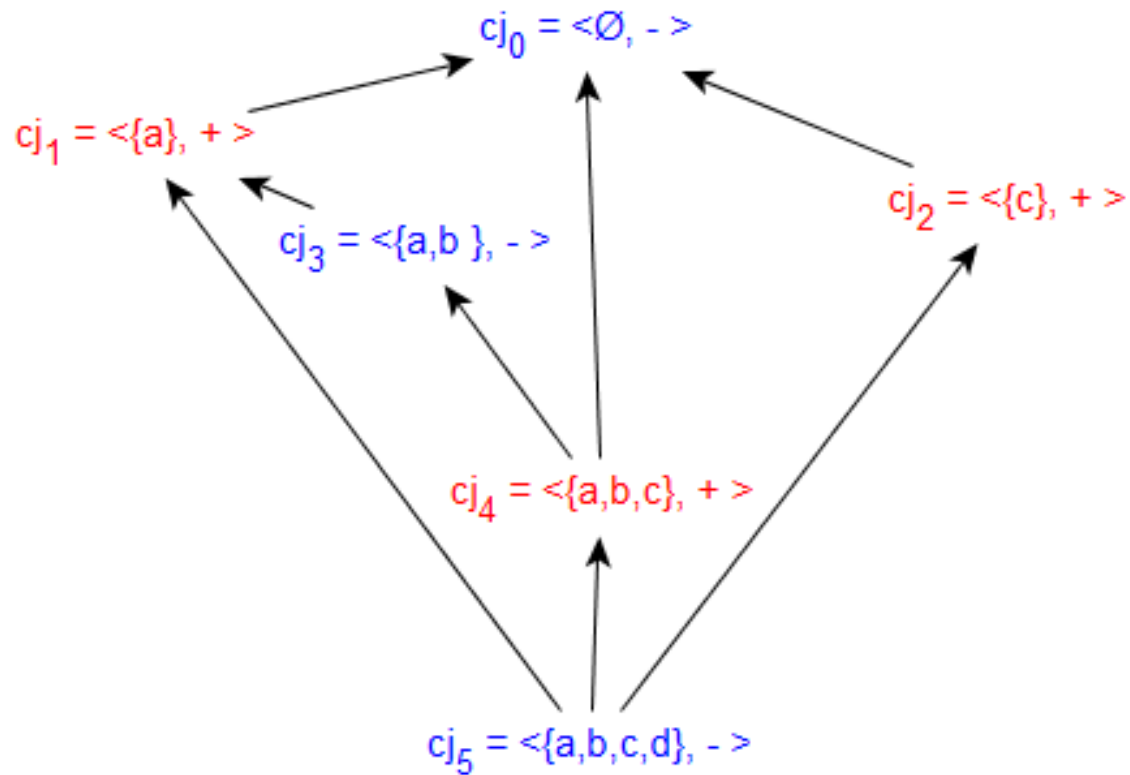
$$F = \{a, b, c, d, e, f\}$$

$$CB = \{ cj_0 = \langle \emptyset, - \rangle, \quad cj_1 = \langle \{a\}, + \rangle, \quad cj_2 = \langle \{c\}, + \rangle, \\ cj_3 = \langle \{a, b\}, - \rangle, \quad cj_4 = \langle \{a, b, c\}, + \rangle, \quad cj_5 = \langle \{a, b, c, d\}, - \rangle \}$$

$$RA = \{ cj_1 \rightarrow_r cj_0, \quad cj_2 \rightarrow_r cj_0, \quad cj_4 \rightarrow_r cj_0, \quad cj_3 \rightarrow_r cj_1, \quad cj_5 \rightarrow_r cj_1, \\ cj_4 \rightarrow_r cj_3, \quad cj_5 \rightarrow_r cj_2, \quad cj_5 \rightarrow_r cj_4 \}$$

Example: Raw Attack Relation

$RA = \{ \langle cj_1, cj_2 \rangle \mid case(cj_1) \supset case(cj_2) \text{ and } judgement(cj_1) \neq judgement(cj_2) \}$



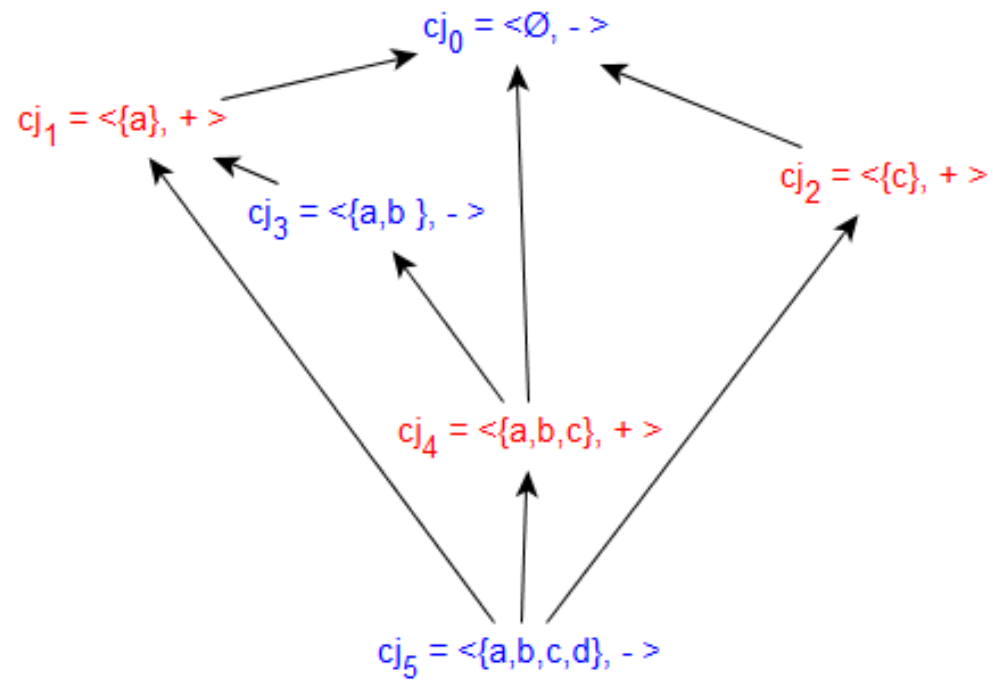
Attack Relation

The **Attack** relation $AT \subseteq RA$ is defined by:

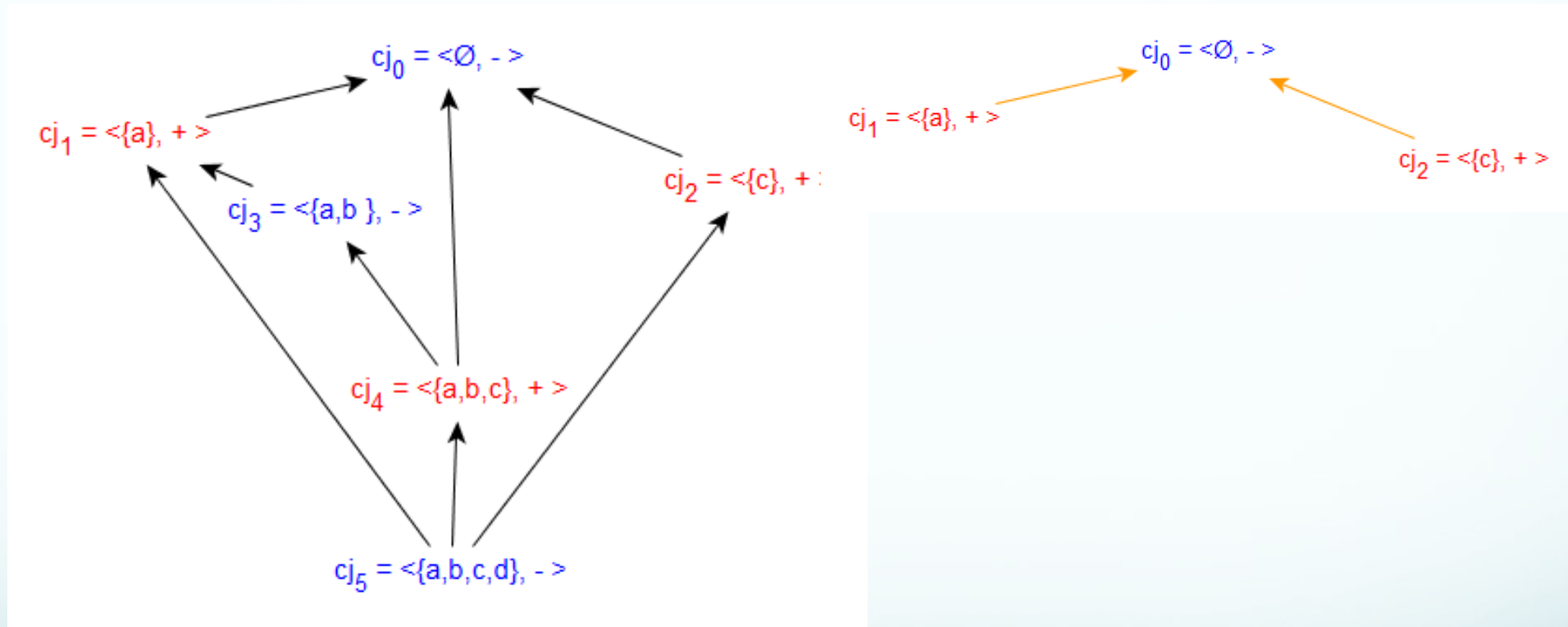
- $\langle cj_1, cj_2 \rangle \in AT$ if $case(cj_2) = \emptyset$ and there is no $cj_3 \rightarrow_r cj_2 \in RA$ such that $case(cj_1) \supset case(cj_3)$
- $\langle cj_1, cj_2 \rangle \in AT$ if there exists $\langle cj_2, cj_4 \rangle \in AT$ and no $cj_5 \rightarrow_r cj_2 \in RA$ such that $case(cj_1) \supset case(cj_5)$
- nothing else is in AT

cj_1 **attacks** cj_2 is denoted $cj_1 \rightarrow cj_2$

Example: Attack Relation

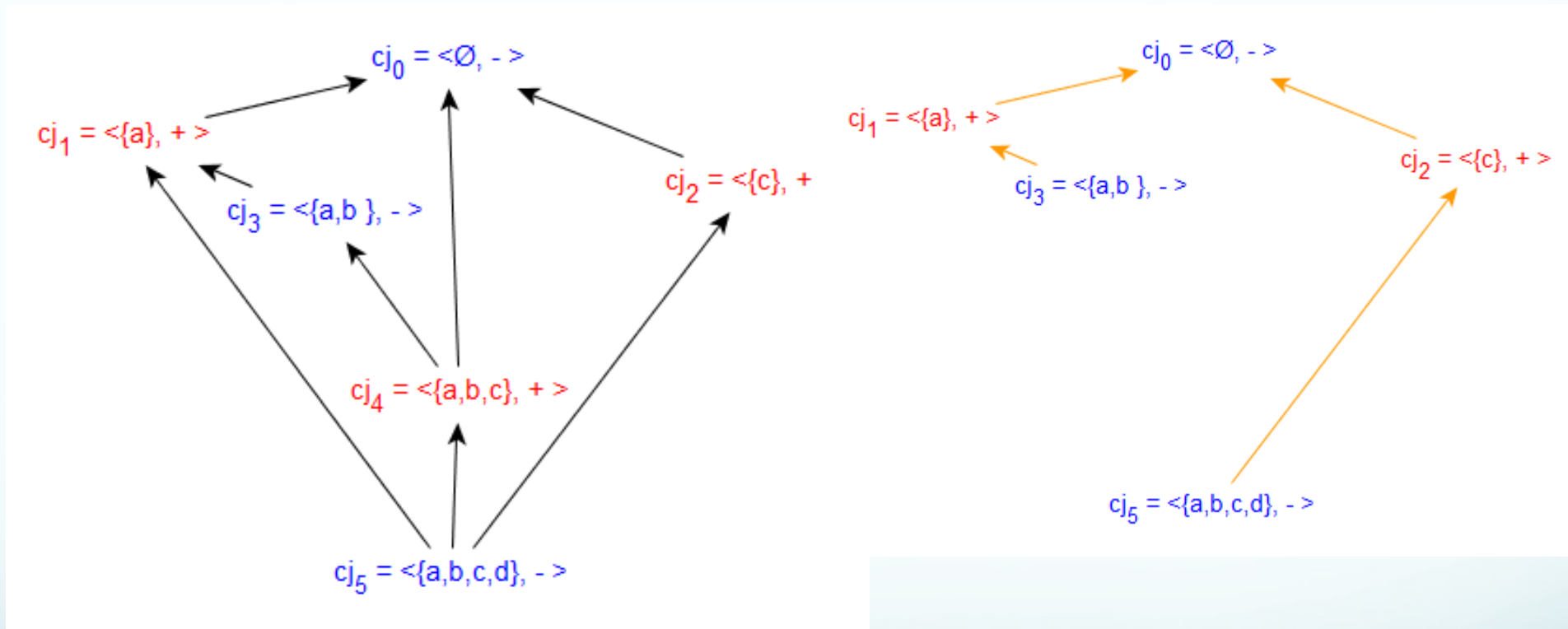
 $cj_0 = \langle \emptyset, - \rangle$

Example: Attack Relation



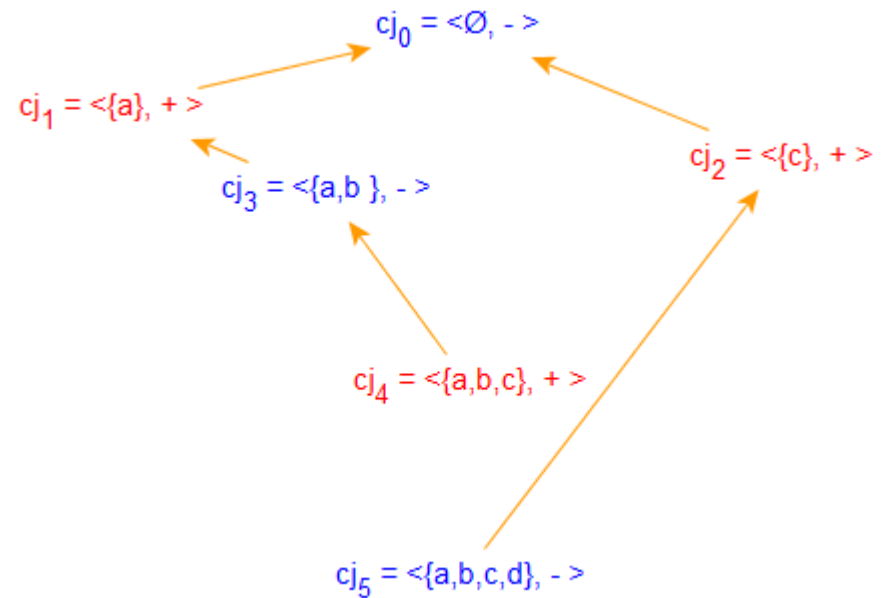
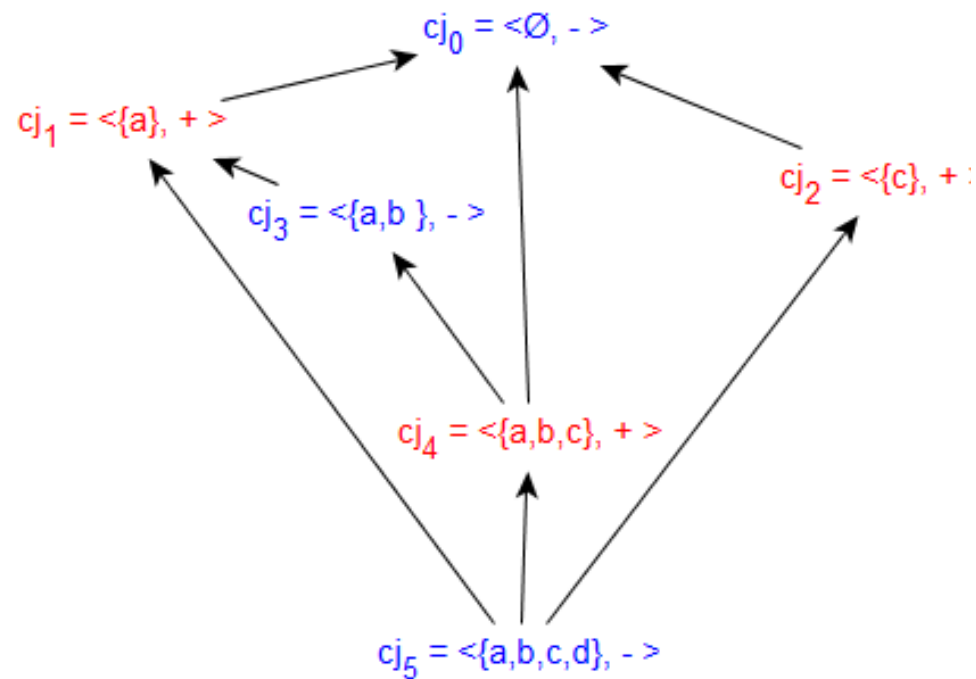
$cj_4 \rightarrow cj_0$ since $case(cj_4) = \{a,b,c\}$ is not minimal factors to overturn the judgement of cj_0

Example: Attack Relation

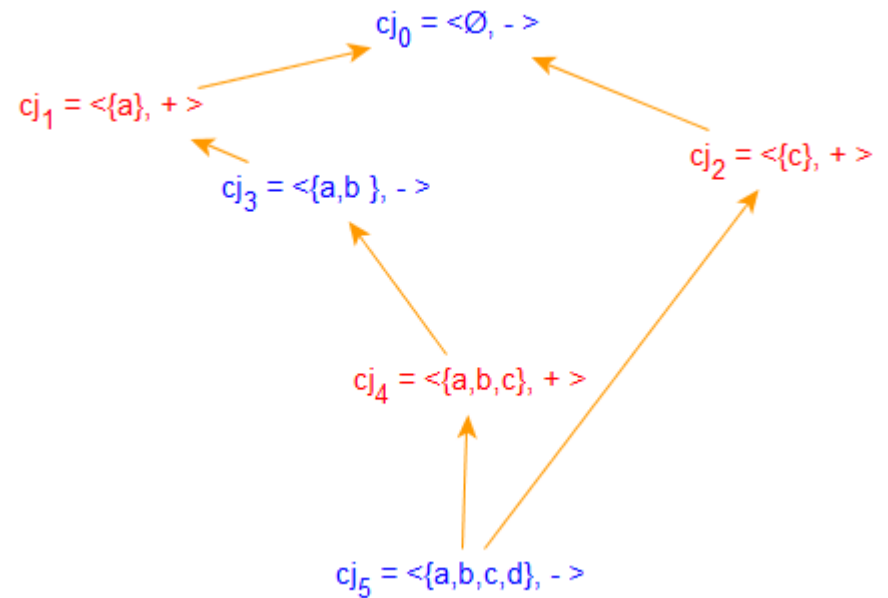
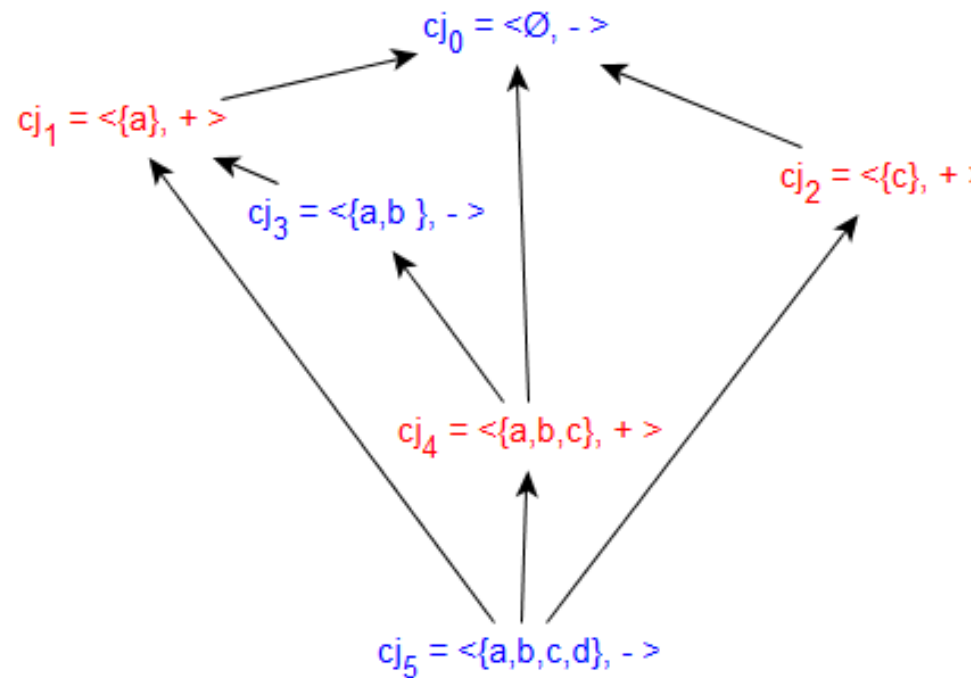


$cj_5 \rightarrow cj_1$ since there is an intermediate overturning a judgement between cj_5 and cj_1

Example: Attack Relation



Example: Attack Relation



Argument

The set of factors responsible for overturning a judgement is called *an argument*.

For each pair $\langle cj_1, cj_2 \rangle$ in the set of attacks AT
 $\alpha(cj_1, cj_2) = case(cj_1) - case(cj_2)$
 is the argument of the attack from cj_1 to cj_2

CB = { $cj_0 = \langle \emptyset, - \rangle$, $cj_1 = \langle \{a\}, + \rangle$, $cj_2 = \langle \{c\}, + \rangle$, $cj_3 = \langle \{a,b\}, - \rangle$,
 $cj_4 = \langle \{a,b,c\}, + \rangle$, $cj_5 = \langle \{a,b,c,d\}, - \rangle$ }

AT = { $cj_1 \rightarrow cj_0$, $cj_2 \rightarrow cj_0$, $cj_3 \rightarrow cj_1$, $cj_4 \rightarrow cj_3$, $cj_5 \rightarrow cj_2$, $cj_5 \rightarrow cj_4$ }

$\alpha(cj_1, cj_0) = \{a\}$, $\alpha(cj_2, cj_0) = \{c\}$, $\alpha(cj_3, cj_1) = \{b\}$,
 $\alpha(cj_4, cj_3) = \{c\}$, $\alpha(cj_5, cj_2) = \{a,b,d\}$, $\alpha(cj_5, cj_4) = \{d\}$

Active Case with Judgement

To predict the judgement of a new case, we look at judgements of similar past cases that have not been overturned.

Given CB , corresponding AT , and a new case c ,
A $c_j \in CB$ is *active* w.r.t. c iff $case(c_j) \subseteq c$ and for all
 $\langle c_{j_n}, c_j \rangle \in AT$, either c_{j_n} is not active w.r.t. c or $case(c_j) \not\subseteq c$.

$$CB = \{ c_{j_0} = \langle \emptyset, - \rangle, c_{j_1} = \langle \{a\}, + \rangle, c_{j_2} = \langle \{c\}, + \rangle, c_{j_3} = \langle \{a,b\}, - \rangle, \\ c_{j_4} = \langle \{a,b,c\}, + \rangle, c_{j_5} = \langle \{a,b,c,d\}, - \rangle \}$$
$$AT = \{ c_{j_1} \rightarrow c_{j_0}, c_{j_2} \rightarrow c_{j_0}, c_{j_3} \rightarrow c_{j_1}, c_{j_4} \rightarrow c_{j_3}, c_{j_5} \rightarrow c_{j_2}, c_{j_5} \rightarrow c_{j_4} \}$$

Let $c_{new} = \{a,b,d\}$ be a new case.

The **active cases with judgment** w.r.t. c_{new} are c_{j_3}, c_{j_0} .

Note: c_{j_1} is not active w.r.t. c_{new} since $\langle c_{j_3}, c_{j_1} \rangle \in AT$ and $case(c_{j_3}) \subseteq c$

Predicted Judgement

Given a CB and corresponding AT and unseen case c ,
The **unique predicted judgement** of c is defined as

$$pj(c) = \text{default } j_0 \text{ iff } \langle \emptyset, j_0 \rangle \text{ is active w.r.t. } c$$

$$CB = \{ cj_0 = \langle \emptyset, - \rangle, cj_1 = \langle \{a\}, + \rangle, cj_2 = \langle \{c\}, + \rangle, cj_3 = \langle \{a,b\}, - \rangle, \\ cj_4 = \langle \{a,b,c\}, + \rangle, cj_5 = \langle \{a,b,c,d\}, - \rangle \}$$

$$AT = \{ cj_1 \rightarrow cj_0, cj_2 \rightarrow cj_0, cj_3 \rightarrow cj_1, cj_4 \rightarrow cj_3, cj_5 \rightarrow cj_2, cj_5 \rightarrow cj_4 \}$$

For case $c_{\text{new}} = \{a,b,d\}$, cj_0 is active w.r.t. c_{new}

$$pj(c) = \text{default judgement}$$

Judgement Theory

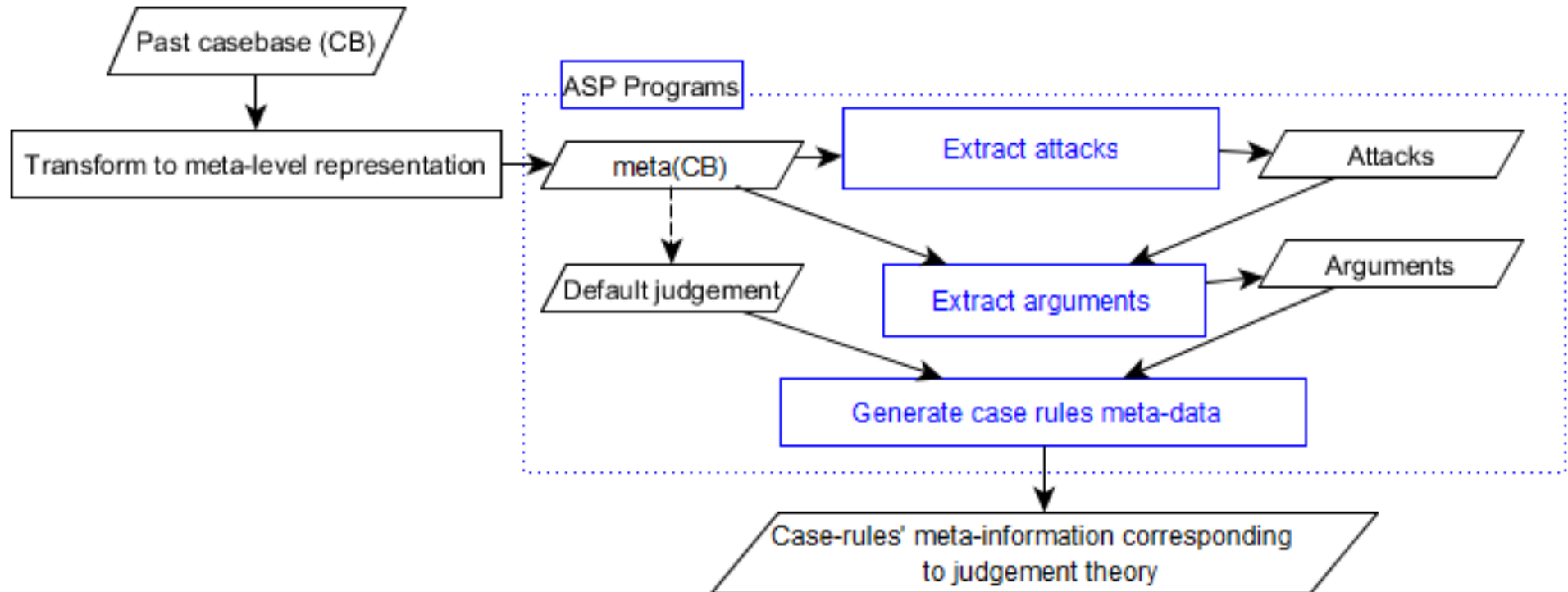
Our aim is to generate a judgement theory from given CB and default j_0 to enable us to infer judgements of new cases c .

For the working example the judgement theory is:

```
def_j0 :- not ab0, not ab1.  
ab0 :- a, not ab2.  
ab1 :- c, not ab5.  
ab2 :- b, not ab3.  
ab3 :- c, not ab4.  
ab4 :- d.
```

The predicted judgement of $c_{\text{new}} = \{a,b,d\}$ is j_0 .

ASP Workflow



The judgement theory is computed using 3 ASP programs

Example



	Factors	Factors	Judgement
cj_0			$\neg dwe$
cj_1	dot		$\neg dwe$
cj_2	dot, ooo		dwe
cj_3	dot, ooo	rpl, far, ria	$\neg dwe$

Past cases with judgement can be seen as “defeasible” rules:

- $\neg dwe$.
- $\neg dwe \leftarrow \text{dot}$.
- $dwe \leftarrow \text{dot, ooo}$.
- $\neg dwe \leftarrow \text{dot, ooo, rpl, far, ria}$.

Each rule can be encoded in a program as a set of meta-level information. For example, for $\neg dwe \leftarrow \text{dot, ooo, rpl, far, ria}$:

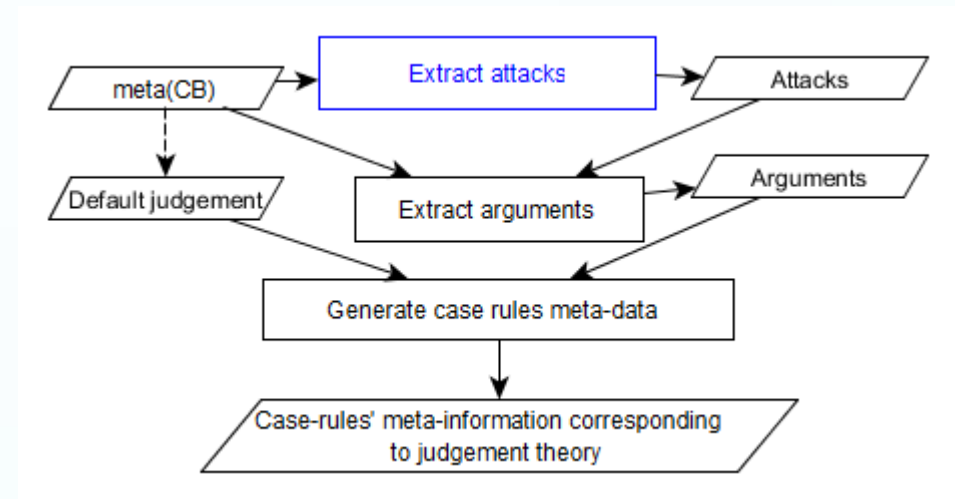
```
cb_id(cj3).
is_rule(cj3,neg_dwe).
in_rule(cj3,neg_dwe,dot).
in_rule(cj3,neg_dwe,ooo).
```

```
in_rule(cj3,neg_dwe,rpl).
in_rule(cj3,neg_dwe,far).
in_rule(cj3,neg_dwe,ria).
```


Example



	Factors	Factors	Judgement
cj_0			$\neg dwe$
cj_1	dot		$\neg dwe$
cj_2	dot, ooo		dwe
cj_3	dot, ooo	rpl, far, ria	$\neg dwe$



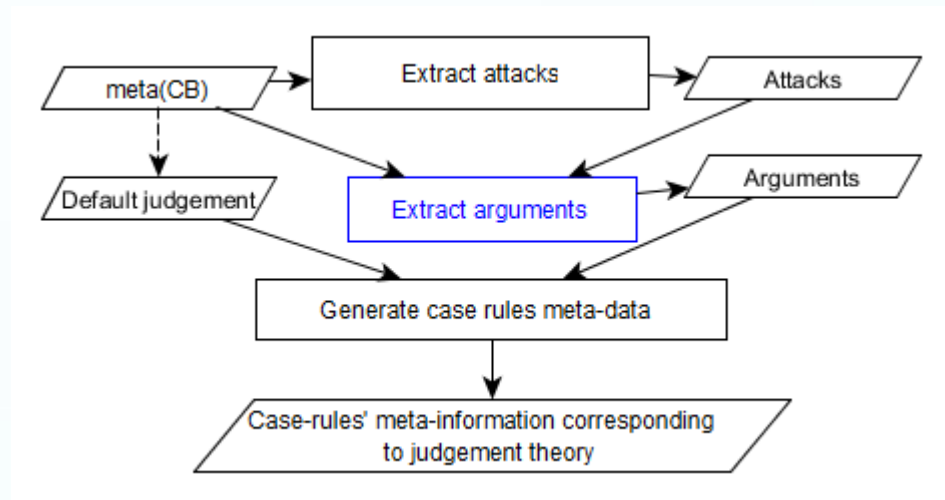
The first program use the casebase to compute the answer set containing all attacks and raw attacks:

$\{ \text{raw_attack}(cj2,c0), \text{raw_attack}(cj3,c2), \text{attack}(cj3,cj2) \}$
 $\text{raw_attack}(cj2,cj1), \text{attack}(cj2,cj0),$

Example



	Factors	Factors	Judgement
cj_0			$\neg dwe$
cj_1	dot		$\neg dwe$
cj_2	dot, ooo		dwe
cj_3	dot, ooo	rpl, far, ria	$\neg dwe$



The second program finds the answer set with all arguments:

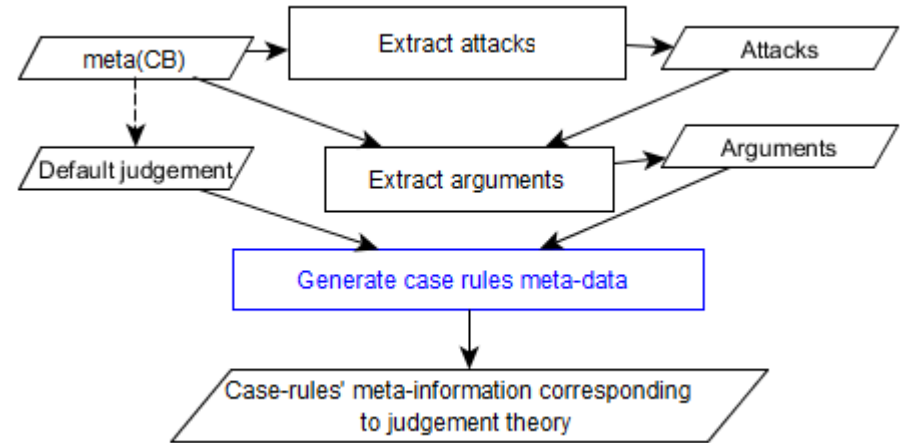
{ argument(cj_2, cj_0, dot),
 argument(cj_3, cj_2, rpl),
 argument(cj_3, cj_2, ria) }

argument(cj_2, cj_0, ooo),
 argument(cj_3, cj_2, far),

Example



	Factors	Factors	Judgement
cj_0			$\neg dwe$
cj_1	dot		$\neg dwe$
cj_2	dot, ooo		dwe
cj_3	dot, ooo	rpl, far, ria	$\neg dwe$



The final program generates the meta-level representation of the case rules:

```

{ is_rule(r0,neg_dwe), in_rule(r0,neg_dwe,not_ab0),
  in_rule(r1,ab0), in_rule(r1,ab0,dot),
  in_rule(r1,ab0,ooo), in_rule(r1,ab0,not_ab1),
  is_rule(r2,ab1), in_rule(r2,ab1,rpl),
  in_rule(r2,ab1,far), in_rule(r2,ab1,ria) }
  
```

Example

The meta-level representation of the judgement theory:

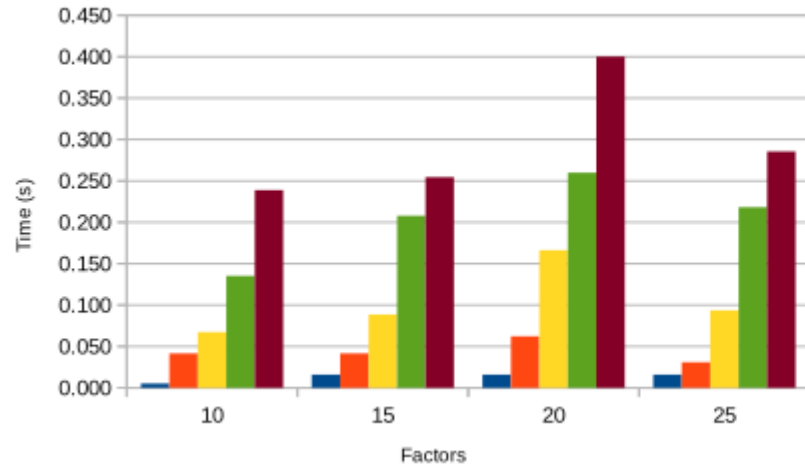
```
{ is_rule(r0,neg_dwe),  in_rule(r0,neg_dwe,not_ab0),  
  in_rule(r1,ab0),      in_rule(r1,ab0,dot),  
  in_rule(r1,ab0,ooo),  in_rule(r1,ab0,not_ab1),  
  is_rule(r2,ab1),      in_rule(r2,ab1,rpl),  
  in_rule(r2,ab1,far),  in_rule(r2,ab1,ria) }
```

Corresponds to the judgement theory:

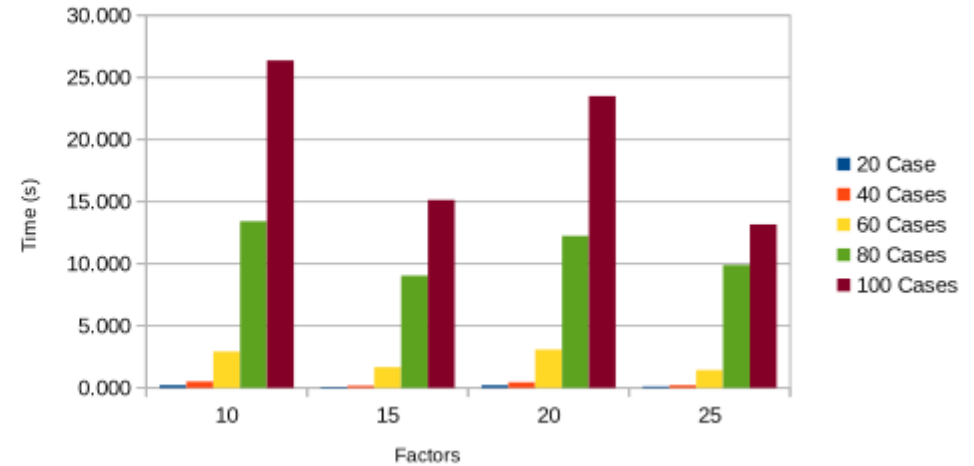
```
¬dwe ← not ab0.  
ab0  ← dot, ooo, not ab1.  
ab1  ← rpl, far, ria.
```

The predicted judgement of {dot, ooo, rpl, far} is dwe.

Evaluation



Time taken to generate the attacks and arguments



Time taken to generate meta-level representation of judgement theory

Correctness Theorem

Given a casebase CB , and associated judgement theory T , and a new case c . Let $A = AS(T \cup c)$.

$$j_0 \in A \text{ if and only if } pj(c) = j_0.$$

Conclusion and Future Work

- We presented a method for reasoning about, and extracting information from, past cases in a casebase to infer the arguments and attacks present in the casebase
- We defined notion of attack to identify factors relevant to judgement of cases and used ASP to infer rules for modelling judgements that can be applied to new cases.

For future work we would like to consider

- how to revise an existing judgement theory
- how to deal with inconsistency among cases
- other possible representations of a casebase