# Esercizi Logica Proposizionale

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Wolfgang Faber Exercises Propositional Logic

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# **Understanding Implication**

#### P ightarrow Q

#### What can we say when P is false?

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# **Understanding Implication**

### P ightarrow Q

Assume *P* represents "has libretto" and *Q* represents "is student".

"If one has a libretto, (s)he is a student."

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# Understanding Implication

Р	Q	$P \rightarrow Q$
"doesn't have libretto"	"is no student"	OK!
"has libretto"	"is no student"	NO!
"doesn't have libretto"	"is student"	OK!
"has libretto"	"is student"	OK!

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# **Understanding Implication**

# Exercise: Which formula represents "Students are exactly those who have a libretto?"

Р	Q	P?Q
"doesn't have libretto"	"is no student"	OK!
"has libretto"	"is no student"	NO!
"doesn't have libretto"	"is student"	NO!
"has libretto"	"is student"	OK!

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## **Understanding Implication**

Exercise: Which formula represents "Students are exactly those who have a libretto?"

Р	Q	<i>P</i> ? <i>Q</i>
"doesn't have libretto"	"is no student"	OK!
"has libretto"	"is no student"	NO!
"doesn't have libretto"	"is student"	NO!
"has libretto"	"is student"	OK!

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### **Eliminate Parentheses**

- Ex. 1.1 from "Logica a Informatica":
  - $\bigcirc ((A \land B) \to (\neg C))$
  - $(A \rightarrow (B \rightarrow (\neg C)))$
  - $((A \land B) \lor (C \to C))$
  - $(\neg (A \lor ((\neg B) \to C)))$
  - $(A \to (B \lor (C \to D)))$

  - $(A \to (B \land ((\neg C) \lor D)))$

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# Tautologies, Contradictions

Ex. 1.3 from "Logica a Informatica": Decide whether the following formulas are tautologies or contradictions:

1 
$$(A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C))$$
  
2  $\neg (A \rightarrow \neg A)$   
3  $A \lor \neg A$   
4  $\bot \rightarrow A$   
5  $\neg A \rightarrow (A \rightarrow B)$   
5  $(A \land B) \land (\neg B \lor C)$   
7  $A \lor B \rightarrow A \land B$   
8  $(A \rightarrow C) \rightarrow ((B \rightarrow C) \rightarrow (A \lor B \rightarrow C))$   
9  $(A \rightarrow B) \rightarrow ((B \rightarrow \neg C) \rightarrow \neg A)$ 

Which of these formulas are satisfiable?

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# Tautologies, Contradictions

Ex. 1.3 from "Logica a Informatica": Decide whether the following formulas are tautologies or contradictions:

$$(A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C))$$

$$(A \rightarrow \neg A)$$

$$(A \rightarrow \neg A)$$

$$(A \rightarrow \neg A)$$

$$(A \rightarrow A)$$

$$(A \rightarrow A)$$

$$(A \rightarrow B) \rightarrow (A \rightarrow B)$$

$$(A \wedge B) \wedge (\neg B \lor C)$$

$$(A \wedge B) \wedge (\neg B \lor C)$$

$$(A \rightarrow B \rightarrow A \land B)$$

$$(A \rightarrow C) \rightarrow ((B \rightarrow C) \rightarrow (A \lor B \rightarrow C))$$

$$(A \rightarrow B) \rightarrow ((B \rightarrow \neg C) \rightarrow \neg A)$$
Which of these formulas are satisfiable?

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## Tautologies, Contradictions

# Similar to Ex. 1.4 from "Logica a Informatica": Decide whether the following formula is satisfiable

$$(A_1 \lor A_2) \land (\neg A_2 \lor \neg A_3) \land (A_3 \lor A_4) \land (\neg A_4 \lor A_5)$$

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# Equivalence, Consequence

- Ex. 1.8 from "Logica a Informatica": Prove that
  - $\bigcirc \bot \lor B \equiv B$
  - **2** $\neg \bot \land B \equiv B$
  - ③ A ⊨ A
  - $A \models B$  and  $B \models C$  implies  $A \models C$

  - **(**)  $\models$  *A* implies  $A \land B \equiv B$
  - $\bigcirc \models A \text{ implies } \neg A \lor B \equiv B$
  - **1** If  $A \models B$  and  $A \models \neg B$  then  $\models \neg A$
  - **(a)** If  $A \models C$  and  $B \models C$  then  $A \lor B \models C$

What are A, B and C?

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## Equivalence, Consequence

Ex. 1.9 from "Logica a Informatica": Check whether

**1** If 
$$A \models B$$
 then  $\neg A \models \neg B$ 

3 If 
$$A \models B$$
 and  $A \land B \models C$  then  $A \models C$ 

If 
$$A \lor B \models A \land B$$
 then  $A \equiv B$ 

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## Transform to equivalent formula in CNF

Ex. 1.13 from "Logica a Informatica": Find equivalent formulas in CNF for

 $(A \rightarrow B) \rightarrow (B \rightarrow \neg C)$  $\neg (A \rightarrow (B \rightarrow \neg C)) \land D$  $\neg (A \land B \land (C \rightarrow D))$  $\neg (A \leftrightarrow B)$ 

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## Transform to equivalent formula in DNF

Ex. 1.13 from "Logica a Informatica": Find equivalent formulas in DNF for

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# Find the Formula!

#### Similar to Ex. 1.10 from "Logica a Informatica": Find f such that

A	В	f
0	0	1
0	1	1
1	0	0
1	1	0

Using only  $\rightarrow$  and  $\perp$ ?

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# Find the Formula!

#### Similar to Ex. 1.10 from "Logica a Informatica": Find f such that

A	В	f
0	0	1
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Using only  $\rightarrow$  and  $\perp$ ?

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# Similar to Ex. 1.13 from "Logica a Informatica": Find f such that

Using only  $\lor$  and  $\neg$ ?

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# Find the Formula!

#### Similar to Ex. 1.13 from "Logica a Informatica": Find f such that

A	В	f
0	0	1
0	1	0
1	0	0
1	1	0

Using only  $\lor$  and  $\neg$ ?

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# Find the Formula in CNF and DNF!

Ex. 1.17 from "Logica a Informatica": Find an f (one in CNF, one in DNF) such that

Α	В	C	f
0	0	0	1
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

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# **Modelling Dinner**

Model the following dinner constraints:

- Available dishes:
  - Farfalle al salmone
  - Risotto agli asparagi
  - Tagliatelle ai funghi
  - Filetto di manzo
  - Spigola grigliata
  - Trancia di pesce spada
- We can choose white or red wine.
- We must choose exactly one primo and one secondo.
- Do not eat fish after mushrooms.
- Choose white wine if fish is involved.

Write a formula such that its models correspond to admissible dinner choices.

# **Reduction to SAT**

Reformulate the following questions such that they can be decided using a SAT algorithm:

- Is  $(P \lor (\neg P \to Q)) \leftrightarrow (P \lor Q)$  valid?
- 2 Does  $P \rightarrow Q$  follow from  $\neg Q \rightarrow \neg P$ ?
- **3** Is  $P \leftrightarrow Q \land P$  a contradiction?
- Is  $P \leftrightarrow P \lor \bot$  a tautology?

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