

Lecture II - part 2

the knowledge base system IDP

June 7, 2016

Inference of the KBS: progress report

Model expansion and visualisation

Imperative + Declarative Programming (IDP)

Inference of the KBS: progress report

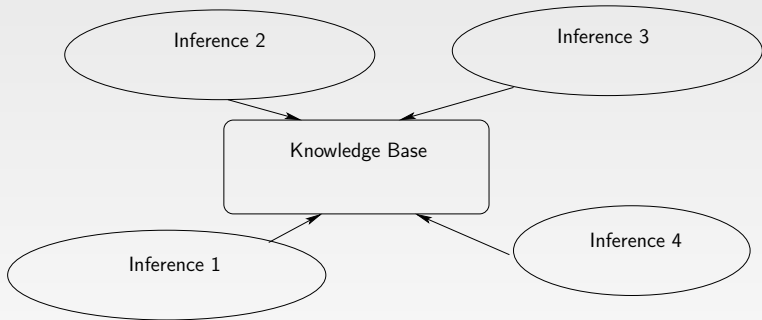
Model expansion and visualisation

Imperative + Declarative Programming (IDP)

How to use Logic for problem solving

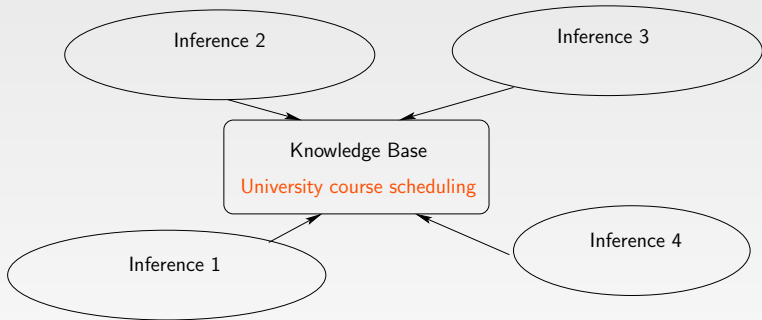
- ▶ A logic theory is a bag of (descriptive) information
 - ▶ A logic theory cannot be executed
 - ▶ A logic theory is not a program
 - ▶ A logic theory is not a representation of a problem
- ▶ So how can we use a logic theory to solve problems?

A Knowledge Base System (KBS)



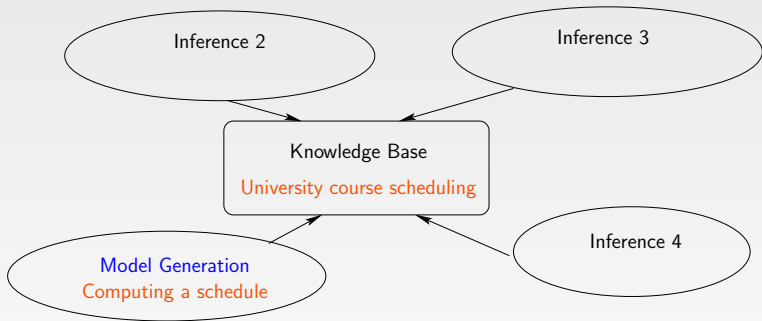
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- ▶ Equipped with different forms of inference:

A Knowledge Base System (KBS)



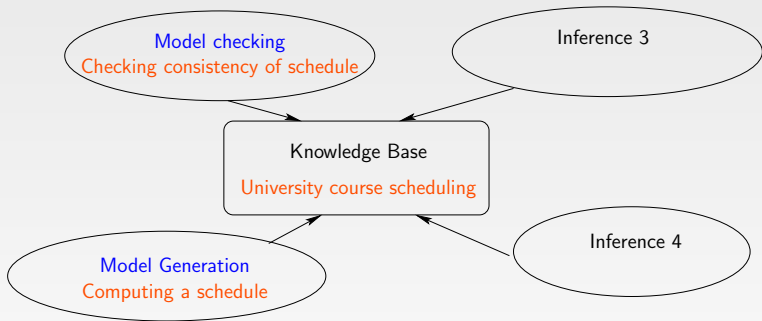
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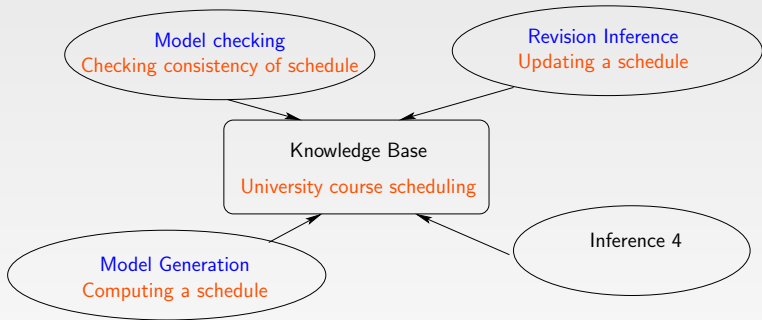
- ▶ Manages a declarative **Knowledge Base (KB)**: a theory
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 - ▶ **Model generation**: Computing a schedule

A Knowledge Base System (KBS)



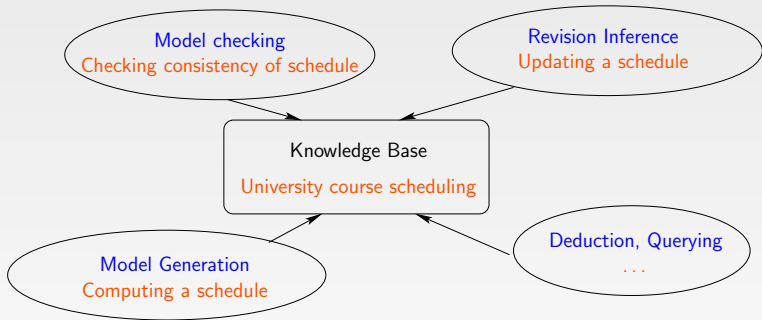
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 - ▶ **Update and Revision**: Updating a given schedule

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- ▶ Manages a declarative **Knowledge Base (KB)**: a theory
- ▶ Equipped with different forms of inference:
 - ▶ **Model generation**: Computing a schedule
 - ▶ **Model checking**: Verifying consistency of a schedule
 - ▶ **Update and Revision**: Updating a given schedule
 - ▶ **Deduction** for verification of the KB
 - ▶ **Querying of defined predicates**, ...

A KBS demo

The course selection demo: interactive configuration

<http://krr.bitbucket.org/courses/>

5 Used forms of inference:

- ▶ Model Checking (P)
- ▶ Propagation (P)
- ▶ Model Generation (NP)
- ▶ Model Generation+Optimization (NP^{NP})
- ▶ Explanation (P)

A KBS demo

Demonstrating a principle that procedural programming languages can't do:

Reusing the same specification/theory/knowledge base to solve different types of problems.

Implementation of KBS

IDP3:

- ▶ A KBS system
- ▶ Programming environment
 - ▶ Programming with theories, structures, inference methods
 - ▶ In an extension of procedural language Lua:

Built by KRR-members Broes De Cat, Bart Bogaerts, Joachim Jansen, Pieter Van Hertum, Jo Devriendt, Ingmar Dasseville (and ex-members Johan Wittocx, Maarten Mariën, Stef De Pooter)

Implementation of KBS

- ▶ Forms of inference currently under development:
 - ▶ (Finite) Model expansion (the core component of IDP3)
 - ▶ Optimisation
 - ▶ Propagation
 - ▶ Querying structures
 - ▶ Δ -model generation and revision:
 - ▶ \sim view materialisation and update in databases
 - ▶ computing & updating defined predicates
 - ▶ Progression of temporal FO(.) theories.
 - ▶ Model revision
 - ▶ Debugging, Explanation.

Model generation/expansion

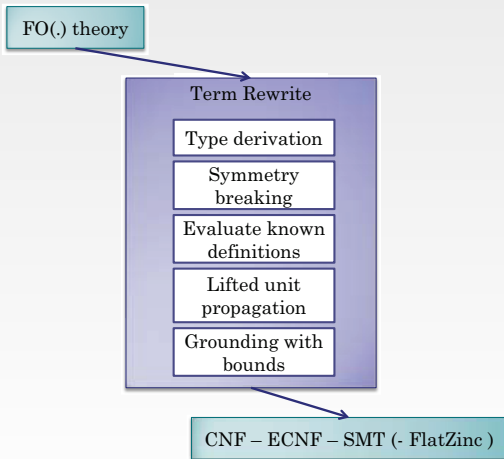
Model Expansion

- ▶ Input:
 - ▶ An FO(.) theory T
 - ▶ An (finite) structure \mathcal{A}_i for a subvocabulary of T , expressing domain and data.
- ▶ Output: a model \mathcal{A} of T expanding \mathcal{A}_i

Special case: Herbrand Model Generation

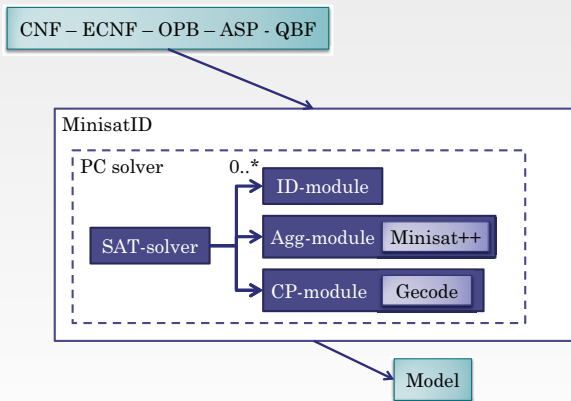
The grounder

Grounding = Eliminating quantification



MinisatID: SMT solver

Solving = computing a model



Technology

Technologies from different computational logic areas integrated:

- ▶ Constraint Programming technology
- ▶ Sat Modulo Theory (SMT)
- ▶ Logic Programming
- ▶ Answer Set Programming
- ▶ MIP: Mixed Integer Programming

Solver	AST (sec.)	PSI (%)
minisatid	950.91	51.62
g12cpx	1126.98	41.68
fzn2smt	1143.47	38.13
ortools	1316.25	30.65
g12lazyfd	1306.10	30.31
gecode	1354.65	29.51
izplus	1350.42	28.05
bprolog	1423.45	24.73
jacop	1435.123	24.67
g12fd	1424.80	23.57
mistral	1525.83	16.91
g12mip	1597.54	12.58

Table : Experimental evaluation of MiniZinc solvers on the CSPs in Benchmark Set B AmadiniGM13.

Benchmark	# solved IDP	# solved Gringo-Clasp
Perm. P. Matching	10	10
Valves Location *	7	4
Still-Life *	2	3
Graceful Graphs	3	9
Bottle Filling	10	10
NoMystery	9	6
Sokoban	7	5
Ricochet Robots	7	10
Crossing Minim. *	0	9
Solitaire	8	9
Weighted Sequence	10	10
Stable Marriage	10	10
Incremental Sched.	6	5
Visit All <i>core</i>	6	7
Knight's Tour <i>core</i>	1	0
Maximal Clique * <i>core</i>	0	1
Graph Colouring <i>core</i>	7	4

Table : Experimental results for benchmarks of the 2013 ASP competition.

Other demos

- ▶ Model expansion with logic-based visualisation
- ▶ Programming with logical inference
- ▶ Temporal Reasoning: planning
- ▶ Temporal reasoning: execution and optimisation

Demos and examples can be accessed from

<https://dtai.cs.kuleuven.be/software/idp/>

IDPd3: logic based visualisation

`dtai.cs.kuleuven.be/krr/idp-ide/?present=SudokuVisualisatie`

- ▶ This application serves to solve sudoku puzzles and to visualise the outcome. The theory T expresses the 3 laws of sudokus: one occurrence of each number in each row, column and block. Model expansion inference solves a puzzle specified in the input structure by returning a model \mathfrak{A} . The solution is *sudoku* ^{\mathfrak{A}} . The user theory T_D3 (see next slide) is mainly a definition of IDPd3 graphical predicate and function symbols defined in terms of symbols interpreted in \mathfrak{A} . Δ -model expansion on T_D3 and input structure \mathfrak{A} computes a value for these predicates which is then transferred to the graphical program *d3*.
- ▶ This is a simple illustration of logic to transform one sort of datastructure in another.

```

theory T_D3 : V_out {
  {
    d3.type(1, Cell(r,k)) = rect <- .
    d3.rect_width(1, Cell(r,k)) = 4 <- .
    d3.rect_height(1, Cell(r,k)) = 4 <- .
    d3_color(1, Cell(r,k)) = "white" <- .
    d3_x(1, Cell(r,k)) = 5*k <- .
    d3_y(1, Cell(r,k)) = 5*r <- .

    d3.type(1, Text(r, k)) = text <- .
    d3_x(1, Text(r, k)) = 5*k <- .
    d3_y(1, Text(r, k)) = 5*r + 1 <- .
    d3.text_size(1, Text(r, k)) = 3 <- .
    d3.text_label(1, Text(r, k)) = t <-
      sudoku(r, k) = c & toString(c) = t.
    d3_color(1, Text(r, k)) = "black" <- .

    d3_order(1, Cell(r, k)) = 0 <- .
    d3_order(1, Text(r,k)) = 1 <- .
  }
}

```

This definition defines slide 1 (argument 1) with:

- ▶ for each cell (r, c) a rectangular object denoted $\text{Cell}(r,c)$, and a text object $\text{Text}(r,k)$.
- ▶ It defines type, width, height, color, x and y position of the rectangular object.
- ▶ It defines type, text size and label, color, x and y position of the text object
- ▶ The **sudoku** number at cell (r, c) is the label of the text object..
- ▶ that text objects are in front of rectangular objects

Δ -model expansion expands a structure interpreting **sudoku** into a structure interpreting all these graphical symbols. This is fed into **d3**.

IDPd3: an example

Input structure

$sudoku = \{1, 1 \rightarrow 1; \dots; 9, 9 \rightarrow 4\}$

...

↓ Δ -model generation

$d3_type = \{1, Cell(1, 1) \rightarrow rect; 1, Text(1, 1) \rightarrow text; \dots\}$

$d3_color = \{1, Cell(1, 1) \rightarrow white; 1, Text(1, 1) \rightarrow black; \dots\}$

$d3_x = \{1, Cell(1, 1) \rightarrow 5; 1, Cell(1, 2) \rightarrow 10; \dots\}$

...

↓ translation to d3 input + d3

A prototype of a knowledge-based programming environment

- ▶ IDP3: A programming environment
- ▶ High level objects: vocabularies, theories, structures
- ▶ Functionalities for manipulation and inference
- ▶ Implemented in the language Lua
- ▶ A new way of mixing Declarative and Procedural knowledge

A demo: generating Sudoku-puzzles

Sudoku-puzzle requirements'

- ▶ (consistency) It should allow one unique solution
- ▶ (minimality) If we delete any value of the puzzle, it has at least two solutions.

Background knowledge base in IDP

Vocabulary

```
vocabulary sudokuVoc {  
  extern vocabulary grid::simpleGridVoc  
  type Num isa nat  
  type Block isa nat  
  Sudoku(Row,Col) : Num  
  InBlock(Block,Row,Col)  
}
```

Theory

```
theory sudokuTheory : sudokuVoc {  
  ! r n : ?1 c : Sudoku(r,c) = n.  
  ! c n : ?1 r : Sudoku(r,c) = n.  
  ! b n : ?1 r c : InBlock(b,r,c) & Sudoku(r,c) = n.  
  ! b r c : InBlock(b,r,c)  
    <=> b = ((r-1)/3)*3 + ((c-1)/3) + 1.  
}
```

A demo: generating Sudoku-puzzles

Puzzle := empty

Generate at most 2 solutions for Puzzle

While 2 solutions were found do{

 Select a random position where the two solutions differ

 Extend Puzzle with the value of the first solution at this position

 Generate at most 2 solutions for Puzzle

}

For each position of Puzzle that contains a value do {

 Delete the value at this position

 Generate at most 2 solutions for Puzzle

 If there are two solutions, undo the deletion of the value.

}

Visualize the puzzle and its unique solution

Procedures

```
procedure createSudoku() {  
    math.randomseed(os.time())  
    local puzzle = grid::makeEmptyGrid(9)  
  
    stdoptions.nrmodels = 2  
    local currssols = modelExpand(sudokuTheory,puzzle)  
  
    while #currssols > 1 do  
        repeat  
            col = math.random(1,9)  
            row = math.random(1,9)  
            num = currssols[1][sudokuVoc::Sudoku](row,col)  
            until num ~= currssols[2][sudokuVoc::Sudoku](row,col)  
  
            makeTrue(puzzle[sudokuVoc::Sudoku].graph,{row,col,num})  
            currssols = modelExpand(sudokuTheory,puzzle)  
        end  
  
        printSudoku(puzzle)  
    }  
}
```

Discussion

Two sorts of inferences:

- ▶ generating solutions to puzzles: **model expansion**
- ▶ Visualizing through **Δ -model expansion**
 - ▶ computing a model of a definition Δ
 - ▶ A special case of model expansion
 - ▶ No search
 - ▶ Can be implemented very differently
 - ▶ = **View materialisation** in deductive databases.

Access and manipulation of **structures**.

- ▶ Structures are objects in the environment
- ▶ Puzzle and its solutions are structures
- ▶ Checking and updating values at positions of puzzle

Reasoning on Temporal theories

- ▶ Temporal theory T : Linear Time Calculus
- ▶ Use Model expansion for planning with optimisation
 - ▶ In the lecture I show a little video showing a idpd3-generated video with the optimal plan to remove all gold.
- ▶ Use [Progression](#) for interactive execution.
 - ▶ Input: T , structure \mathcal{A} representing state at time i
 - ▶ Output: structure \mathcal{A}' representing possible state at time $i + 1$.
- ▶ Illustration: pacman.

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- ▶ Illustration: pacman.

IDP is the only system that we know of that can use the same formal specification to solve both tasks.

The pacman demo:

- ▶ Go to the IDP-web page
`https://dtai.cs.kuleuven.be/software/idp/`
- ▶ Let column: select "Demos"
- ▶ Next page contains different visualised demos of IDP. Select bottom middle "More visualised demos."
- ▶ Next page, select link below "Pacman".
- ▶ Next page, press Run
- ▶ A labyrinth with above, a green and black rectangles appear, indicating possible move actions van the pacman. Click the black rectangles and pacman moves in the right direction.
- ▶ Due to memory limitations on the server, only a few moves can be played.
- ▶ To avoid memory limitations, download the system.

A company running standard software systems on this principle:

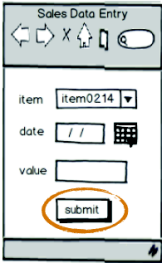
- ▶ LogicBlox - Datalog :
<http://www.logicblox.com/>

File Edit View Document Tools Window Help

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UI event handling



```
^sales[prod,date,store]=value
<-
+button_clicked(form,submit),
sales_entry_form_user(form,user),
dropdown_selected[form]=prod,
date_fld_value[form,date_fld]=date,
num_fld_value[form,val_fld]=value,
manager(store,user).
```

Thursday, March 1, 12

LOGICBLOX®

some clients



Generic interactive configuration

<http://krr.bitbucket.org/autoconfig/>

Another application, selectable from the demos webpage (left middle box) is interactive decision enactment = business logic terminology for generic interactive configuration.

- ▶ Automatic generation of a window with values for all ground literals of a configuration problem
 - ▶ Three-valued structure underneath
- ▶ User can select values
 - ▶ System propagates user choices
 - ▶ Optimal propagation versus approximate propagation
- ▶ System expands current partial structure on demand
- ▶ System searches model minimizing a cost term on demand
- ▶ Generic, incremental:
 - ▶ Adding symbols
 - ▶ Other vocabularies

The application is a benchmark problem of business rules

- ▶ USERV
- ▶ Deciding car insurance policy for client.

To be submitted to RuleML.

Advantages compared to business rule systems

- ▶ Clear declarative semantics: formal and informal
- ▶ Reasoning backward from desired outcome.
- ▶ Increased functionalities
 - ▶ computing solutions under uncertainty
 - ▶ propagation
 - ▶ optimisation
 - ▶ explanation (is not implemented)
 - ▶ ...

Integrating logic in imperative languages

Joost Vennekens: Lowering the learning curve for declarative programming: a Python API for the IDP system. International Workshop on User-Oriented Logic Programming (IULP 2015) 31st August 2015, Cork (Ireland) CoRR abs/1511.00916


```
idp.Type("Number", range(10))
idp.Function("Given(Square):  Number", d)
idp.Function("Sol(Square):  Number")
idp.Define("Diff(Square, Square)", "lambda x,y:  x !=
y and (SameRow(x,y) or SameCol(x,y) or
SameSmallSq(x,y))")
idp.Constraint("all(Sol(x) == Given(x) for x in
Square if Given(x) != 0)")
idp.Constraint("all(Sol(x) != 0 for x in Square)")
idp.Constraint("all(Sol(x) != Sol(y) for (x,y) in
Diff)")
show(idp.Sol)
```

See webpages:

- ▶ <https://dtai.cs.kuleuven.be/software/idp/>
- ▶ Online IDE (edit, save, download, run, many examples)
- ▶ Demos
 - ▶ map coloring
 - ▶ interactive course selection
 - ▶ more visualised demos
 - ▶ pacman - interactive execution
 - ▶ Science Week example : visualisation of errors
- ▶ tutorial, manual
- ▶ ICLP-contest 2015 solutions

Future

Knowledge-based software engineering

- ▶ Important gains to be made:
 - ▶ development time
 - ▶ compactness
 - ▶ correctness
 - ▶ reuse
 - ▶ maintainability
- ▶ Great scientific and practical challenges

We are in the process of searching niches with industry where our technology could already make a difference

A KRR-team with Ingmar Dasseville, Jo Devriendt and Matthias van der Hallen won the International Logic Programming and Constraint Programming competition with IDP.